



minimum number of clock cycles — that the instruction takes; so the actual time taken will also depend on the clock rate. A common clock rate for 6809 systems is one MHz — one million cycles per second. Thus, each clock cycle takes one millionth of a second. The straightforward instruction:

LDA DATIM

using a 16-bit address takes five cycles, and so executes in five millionths of a second. The instruction:

PSHS PC, B, CC

takes five cycles plus one cycle for each byte that is pushed onto the stack; in this case a total of nine cycles (remember that the PC is two bytes).

If a system does not include a real-time clock then the only way to measure elapsed time is by means of a software delay routine. This executes a sequence of instructions whose individual times have been chosen so that the sum gives the required interval. Such intervals are usually measured in milliseconds (thousandths of a second), so there is no need to be too exact — the odd millionth of a second will not matter. Assuming a clock rate of one MHz, the Software Delay routine we give here will produce delays in the range 1 to 255 milliseconds: the exact number of milliseconds (ms) being passed as a parameter in A. The notation (A) means the contents of accumulator A.

The calculation to find the constant COUNT can be expressed as follows:

Instruction	Number Of Clock Cycles	Number Of Times Executed	Time Taken (Clock Cycles)
PSHS B,CC	7	1	7
LDB #COUNT	2	(A)	(A) * 2
DECB	2	(A) * COUNT	(A) * COUNT * 2
BNE LOOP2	3	(A) * COUNT	(A) * COUNT * 3
DECA	2	(A)	(A) * 2
BNE LOOP1	3	(A)	(A) * 3
PULS PC,B,CC	9	1	9

This gives a total of $(A) * (7 + 5 * COUNT) + 16$ clock cycles. To make the calculation easier, we will ignore the 16. At a clock rate of one MHz, there are 1000 clock cycles in a millisecond so the total time should be $(A) * 1000$ clock cycles.

$$(A) * (7 + 5 * COUNT) = (A) * 1000$$

$$(7 + 5 * COUNT) = 1000$$

$$5 * COUNT = 973$$

$$COUNT = 195 \text{ (to the nearest integer)}$$

It is quite feasible to make more accurate delays, and to use the 16-bit registers for a greater range, but the principle of decrementing a register a fixed number of times remains the same.

Terminal Emulation Routine

ESCAPE	EQU	27	
SPACE	EQU	32	(Space is ASCII 32)
OUTCH	EQU		Enter operating system address here
	ORG	\$1000	
CTABLE	RMB	32	Table of control characters
ETABLE	RMB	128	Table of Escape characters
EFLAG	FCB	0	Flag to indicate whether last character was an Escape
DISPCH	PSHS	X	Save X
	TST	EFLAG, PCR	Check if last character was an Escape
	BEQ	DISP1	If not an Escape, then go to DISP1
	LEAX	ETABLE, PCR	Else get address of ETABLE in X
	LDA	A,X	Get replacement character using the original character in A as the offset
	CLR	EFLAG,PCR	Reset EFLAG
	BRA	FINISH	
DISP1	CMPA	SPACE	Check if control character
	BGE	FINISH	If not control character, then goto FINISH
	CMPA	ESCAPE	Else check if Escape
	BEQ	ESCCH	If it is Escape, then goto ESCCH
	LEAX	CTABLE,PCR	Get address of CTABLE in X
	LDA	A,X	Get replacement character using the original character in A as offset
	BRA	FINISH	
ESCCH	INC	EFLAG,PCR	Set EFLAG to indicate character was Escape
FINISH	PULS	X	Restore X
	JMP	OUTCH	Display character in A
	END		Note that the RTS at the end of OUTCH will return control from here to the calling program

Software Delay Routine

COUNT	EQU	195	Subroutine to delay (A) milliseconds
	ORG	\$1000	See calculation
DELAY	PSHS	B,CC	Save the other two registers affected
LOOP1	LDB	#COUNT	Count for 1 ms
LOOP2	DECB		Keep decrementing
	BNE	LOOP2	Until B reaches zero
	DECA		Decrement A after each ms
	BNE	LOOP1	Until A reaches zero
	PULS	PC,B,CC	Return